

On Covering Problems of Rado

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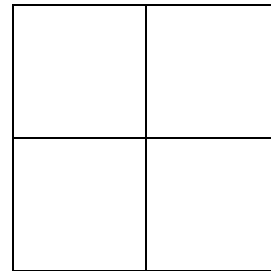
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$$F(S) \leq 1/4$$

A Conjecture by T. Rado in 1928

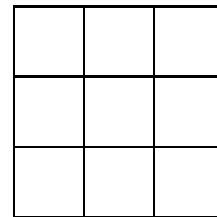
For any set S of axis-parallel squares in the plane, there is always a subset \mathcal{I} of pairwise disjoint squares such that

total area of squares in \mathcal{I}

is at least $1/4$ of

union area of squares in S

Lower Bound



$$F(S) \geq 1/9$$

Greedy algorithm: largest first (T. Rado 1928)

$$F(S) \geq 1/8.6$$

Semi-greedy algorithm: largest or neighbors (Zalgaller 1960)

More Precisely...

The Question

For a finite set S of convex sets in \mathbb{R}^d , the *union volume* (or *union area* when $d = 2$) of S is

$$|S| = |\cup_{C \in S} C|.$$

For a convex set S in \mathbb{R}^d , define

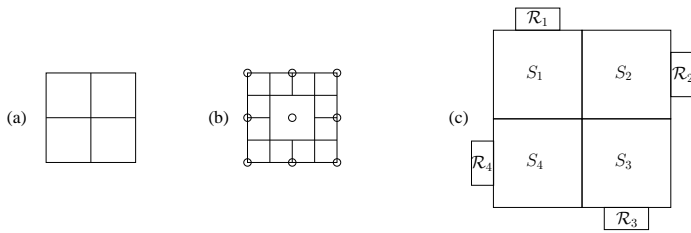
$$F(S) = \inf_S \sup_{\mathcal{I}} \frac{|\mathcal{I}|}{|S|},$$

where S ranges over all finite sets of convex sets in \mathbb{R}^d that are homothetic to S , and \mathcal{I} ranges over all independent subsets of S .

$$F(S) \stackrel{?}{=} 1/4$$

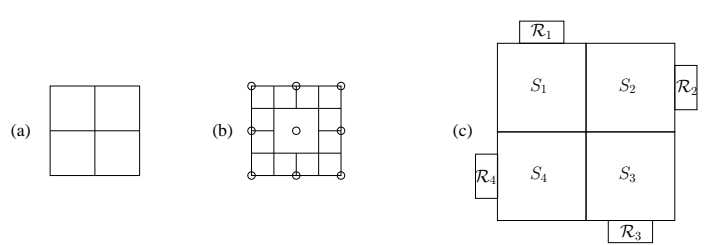
Is it true that $F(S) \geq 1/4$ for a square S ?

Disproved by Ajtai in 1973



- (a) A system of 4 squares.
- (b) An ambiguous system of 13 squares.
- (c) Ajtai's construction shown schematically: $F(S) < 1/4$.

Improve Ajtai's System



- Step 1: design a better gadget \mathcal{R}
- Step 2: repeat the system in a tiling

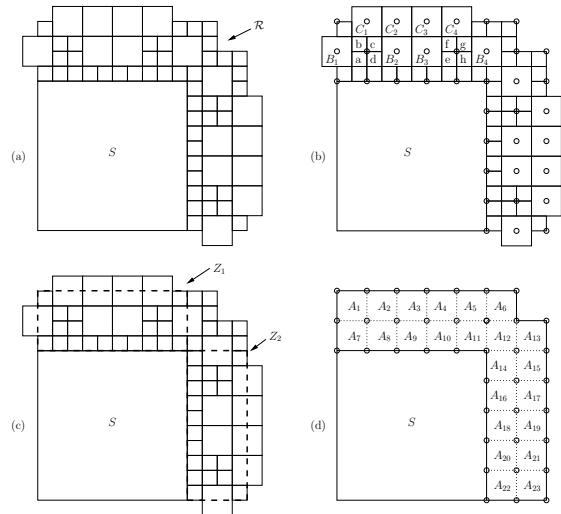
Our Results

For a square S ,

$$1/8.6 < 1/8.4797 \dots \leq F(S) \leq 1/4.0421 \dots < 1/4.$$

Improved $F(S)$ bounds for disks, centrally symmetric convex sets in \mathbb{R}^2 , arbitrary convex sets in \mathbb{R}^2 , and hypercubes in \mathbb{R}^d .

Improved $f(S)$ bounds for *congruent* convex sets in \mathbb{R}^2 .



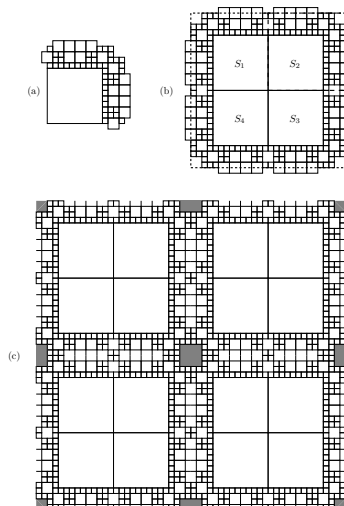
Improve Upper Bound

Estimate of Ajtai 1973:

$$1/4 - 1/1728 = 1/4.0092 \dots$$

Our result:

$$1/4 - 1/384 = 1/4.0421 \dots$$



Improve Lower Bound

T. Rado 1928 (largest first): $1/9$

R. Rado 1949 (largest or two neighbors): $1/8.75$

Zalgaller 1960 (largest or at most four neighbors): $1/8.6$

Our results:

largest or two neighbors

or two neighbors of a neighbor

or two neighbors of a neighbor of a neighbor...

$1/8.5699\dots$

largest or three large neighbors

$1/8.4797\dots$

In Every Iteration of Step 3

Since S_i and S_j are two squares that intersect S_k and touch two opposite sides of T_k ,

$$x_i + x_j + x_k \geq y_k = x_k + z_k.$$

Then,

$$x_i + x_j \geq z_k.$$

So when we reach the last iteration of Step 3...

Algorithm A1

- $\mathcal{S} = \{S_1, \dots, S_n\}$ is a set of axis-parallel squares.
- T_i is the smallest axis-parallel square that contains all squares in \mathcal{S} that intersect S_i .
- x_i is the side length of S_i .
- y_i is the side length of T_i .
- $z_i = y_i - x_i$.
- λ is a constant (will be set to $8.5699\dots$).

To construct an independent set \mathcal{I} , initialize \mathcal{I} to be empty, then repeat the following *selection round* until \mathcal{S} is empty.

$$\begin{aligned} \frac{|T_i \cup T_j|}{|S_i| + |S_j|} &= \frac{|T_i| + |T_j| - |T_i \cap T_j|}{|S_i| + |S_j|} \\ &\leq \frac{(x_i + z_k)^2 + (x_j + z_k)^2}{x_i^2 + x_j^2} - \frac{x_k^2}{x_i^2 + x_j^2} \\ &= 1 + \frac{2z_k(x_i + x_j) + 2z_k^2}{x_i^2 + x_j^2} - \frac{x_k^2}{x_i^2 + x_j^2} \\ &\leq 1 + \frac{2z_k(x_i + x_j) + 2z_k^2}{(x_i + x_j)^2/2} - \frac{(\sqrt{\lambda} - 2)^2}{2} \\ &= 1 + \frac{4z_k}{x_i + x_j} + \frac{4z_k^2}{(x_i + x_j)^2} - \frac{(\sqrt{\lambda} - 2)^2}{2} \\ &\leq 9 - (\sqrt{\lambda} - 2)^2/2 \end{aligned}$$

Selection Round of Algorithm A1

1. Find the largest square S_l in \mathcal{S} .
Assume without loss of generality that $x_l = 1$.
2. If $y_l \leq \sqrt{\lambda}$, add S_l to \mathcal{I} , delete from \mathcal{S} the squares that intersect S_l , then return. Otherwise, set $k \leftarrow l$ and continue with the next step.
3. Let S_i and S_j be two squares that intersect S_k and touch two opposite sides of T_k . If both z_i and z_j are at most z_k , add S_i and S_j to \mathcal{I} , delete from \mathcal{S} the squares that intersect S_i or S_j , then return. Otherwise, set $k \leftarrow i$ or j such that z_k increases, then repeat this step.

The idea: keep searching if a neighbor has "more" neighbors.

Two Cases

1. If S_l is selected (in Step 2),

$$\frac{|T_l|}{|S_l|} \leq \lambda.$$

2. If S_i and S_j are selected (in the last iteration of Step 3),

$$\frac{|T_i \cup T_j|}{|S_i| + |S_j|} \leq 9 - (\sqrt{\lambda} - 2)^2/2.$$

$$9 - (\sqrt{\lambda} - 2)^2/2 = \lambda$$

$$\lambda = (\sqrt{46} + 2)^2/9 = 8.5699\dots$$

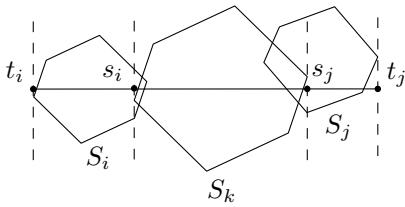
$$F(S) \geq 1/\lambda \text{ for a square } S.$$

More Lower Bounds

Summary

$$3^d - (\lambda_d^{1/d} - 2)^{d/2} = \lambda_d$$

$F(S) \geq 1/\lambda_d$ for a hypercube S in \mathbb{R}^d .



For a square S ,

$$1/8.6 < 1/8.4797\dots \leq F(S) \leq 1/4.0421\dots < 1/4.$$

Improved $F(S)$ bounds for disks, centrally symmetric convex sets in \mathbb{R}^2 , arbitrary convex sets in \mathbb{R}^2 , and hypercubes in \mathbb{R}^d .

Improved $f(S)$ bounds for *congruent* convex sets in \mathbb{R}^2 .

Thank you!

$F(S) \geq 1/8.5699\dots$ for a centrally symmetric convex set in \mathbb{R}^2 .

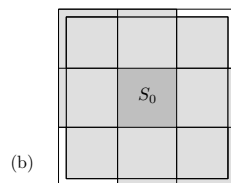
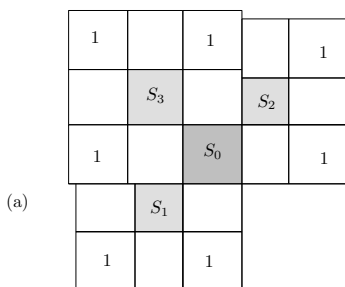
$F(S) \geq 1/8.3539\dots$ for a disk (by a tighter analysis).

Selection Round of Algorithm A2

- Let S_0 be the largest square in \mathcal{S} .
Assume without loss of generality that S_0 is a unit square.
Let $\mathcal{S}_0 \subseteq \mathcal{S} \setminus \{S_0\}$ be the set of squares of side length at least $s = 0.8601\dots$ that intersect S_0 .
- If \mathcal{S}_0 contains three disjoint squares S_1 , S_2 , and S_3 , then add S_1 , S_2 , and S_3 to \mathcal{I} . Otherwise add S_0 to \mathcal{I} .
- For each square S_i added to \mathcal{I} , remove from \mathcal{S} the squares that intersect S_i .

The idea: take the largest or three *large* neighbors.

Two Cases



$$\frac{8 + 3s^2 + 10s}{3s^2} = 7 + 2s^2 = 8.4797\dots$$